

CS250P: Computer Systems Architecture Explicit Parallelism



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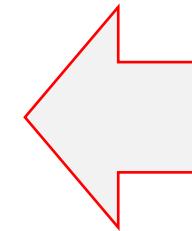
Fall 2023



Large amount of material adapted from MIT 6.004, “Computation Structures”,
Morgan Kaufmann “Computer Organization and Design: The Hardware/Software Interface: RISC-V Edition”,
and CS 152 Slides by Isaac Scherson

Modern Processor Topics - Performance

- ❑ Transparent Performance Improvements
 - Pipelining, Caches
 - Superscalar, Out-of-Order, Branch Prediction, Speculation, ...
 - Covered in CS250A and others
- ❑ Explicit Performance Improvements
 - SIMD extensions, AES extensions, ...
 - ...

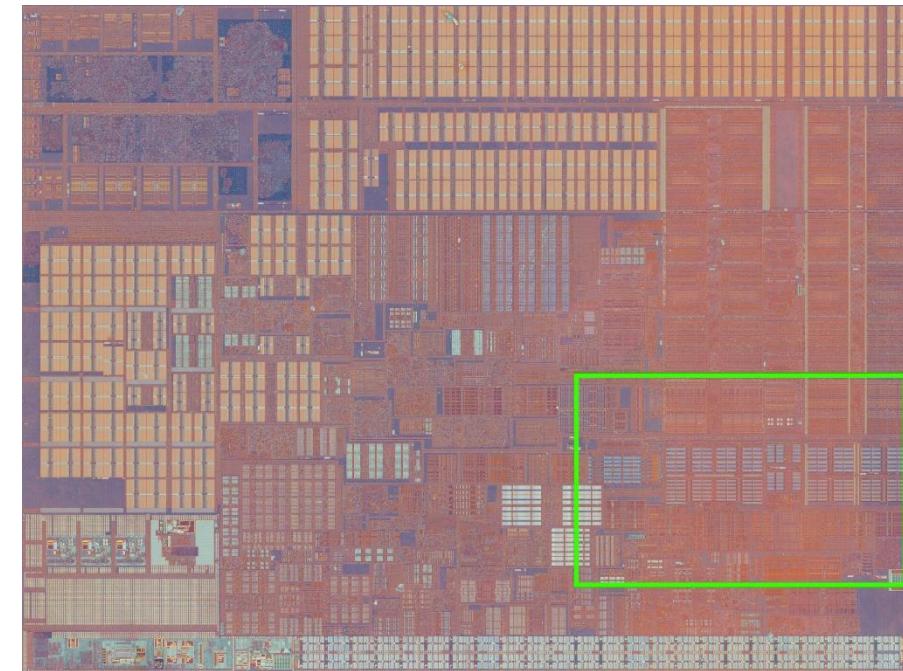


SIMD operations

- ❑ Single ISA instruction performs same computation on multiple data
- ❑ Typically implemented with special, wider registers
- ❑ Example operation:
 - Load 32 bytes from memory to special register X
 - Load 32 bytes from memory to special register Y
 - Perform addition between each 4-byte value in X and each 4 byte value in Y
 - Store the four results in special register Z
For i in (0 to 7): $Z[i] = X[i] + Y[i];$
 - Store Z to memory
- ❑ RISC-V SIMD extensions (P) is still being worked on (as of 2021)

Example: Intel SIMD Extensions

- ❑ More transistors (Moore's law) but no faster clock, no more ILP...
 - More capabilities per processor has to be explicit!
- ❑ New instructions, new registers
 - Must be used explicitly by programmer or compiler!
- ❑ Introduced in phases/groups of functionality
 - SSE – SSE4 (1999 – 2006)
 - 128 bit width operations
 - AVX, FMA, AVX2, AVX-512 (2008 – 2019)
 - 256 – 512 bit width operations
 - F16C, and more to come?

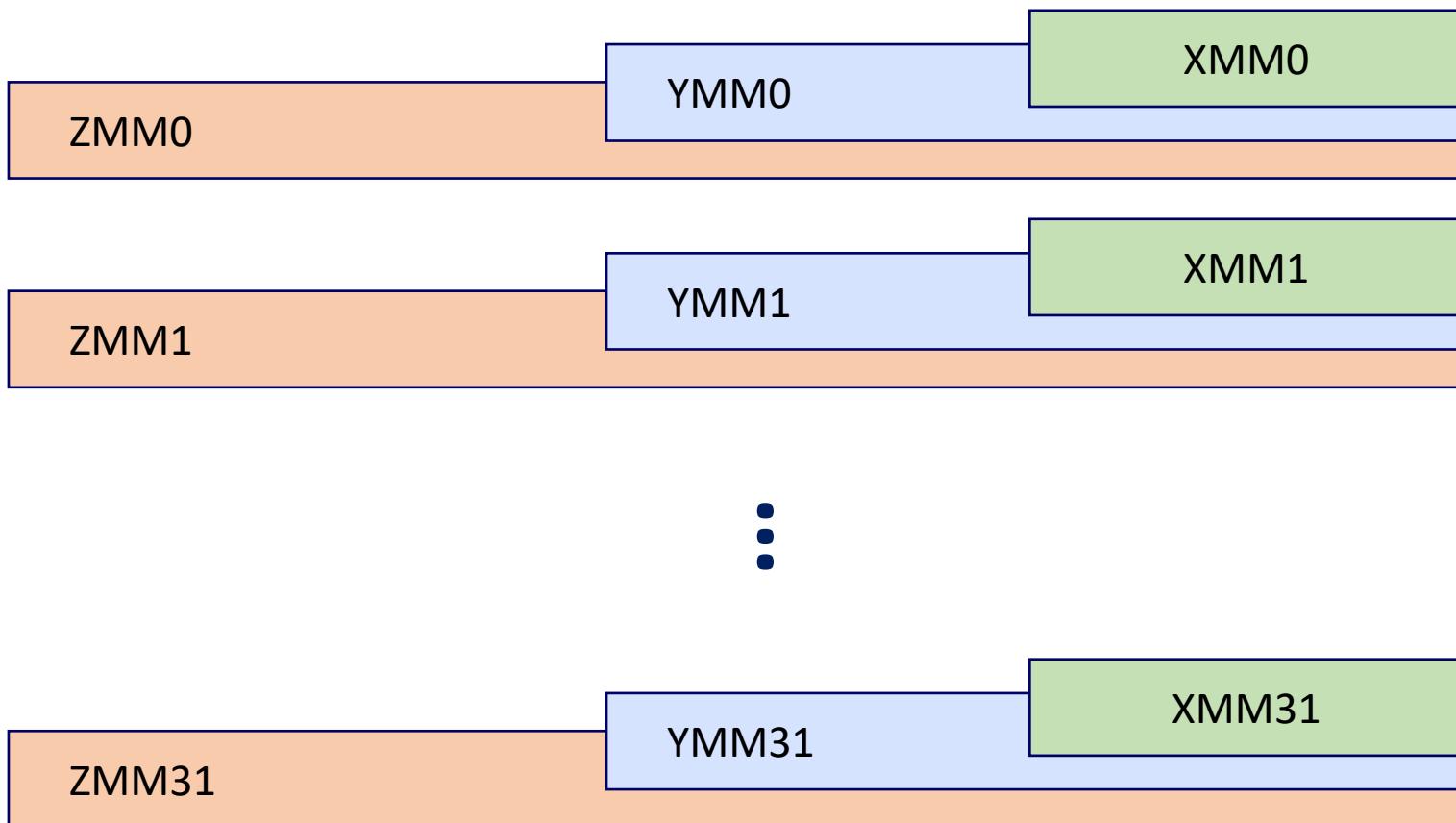


Aside: Do I Have SIMD Capabilities?

❑ less /proc/cpuinfo

```
flags          : fpu vme de pse tsc msr pae mce cx8 apic sep mttr pge mca cmov pat p  
se36 clflush dts acpi mmx fxsr sse sse2 ss ht tm pbe syscall nx pdpe1gb rdtscp lm con  
stant_tsc art arch_perfmon pebs bts rep_good nopl xtopology nonstop_tsc cpuid aperfmp  
erf tsc_known_freq pni pclmulqdq dtes64 monitor ds_cpl vmx est tm2 ssse3 sdbg fma cx1  
6 xtpr pdcm pcid sse4_1 sse4_2 x2apic movbe popcnt tsc_deadline_timer aes xsave avx f  
16c rdrandlahf_lm abm 3dnowprefetch cpuid_fault epb invpcid_single pti ssbd ibrs ibp  
b stibp tpr_shadow vnmi flexpriority ept vpid fsgsbbase tsc_adjust bmi1 avx2 smep bmi2  
erms invpcid mpx rdseed adx smap clflushopt intel_pt xsaveopt xsavewc xgetbv1 xsaves  
dtherm ida arat pln pts hwp hwp_notify hwp_act_window hwp_epp flush_lld
```

Intel SIMD Registers (AVX-512)



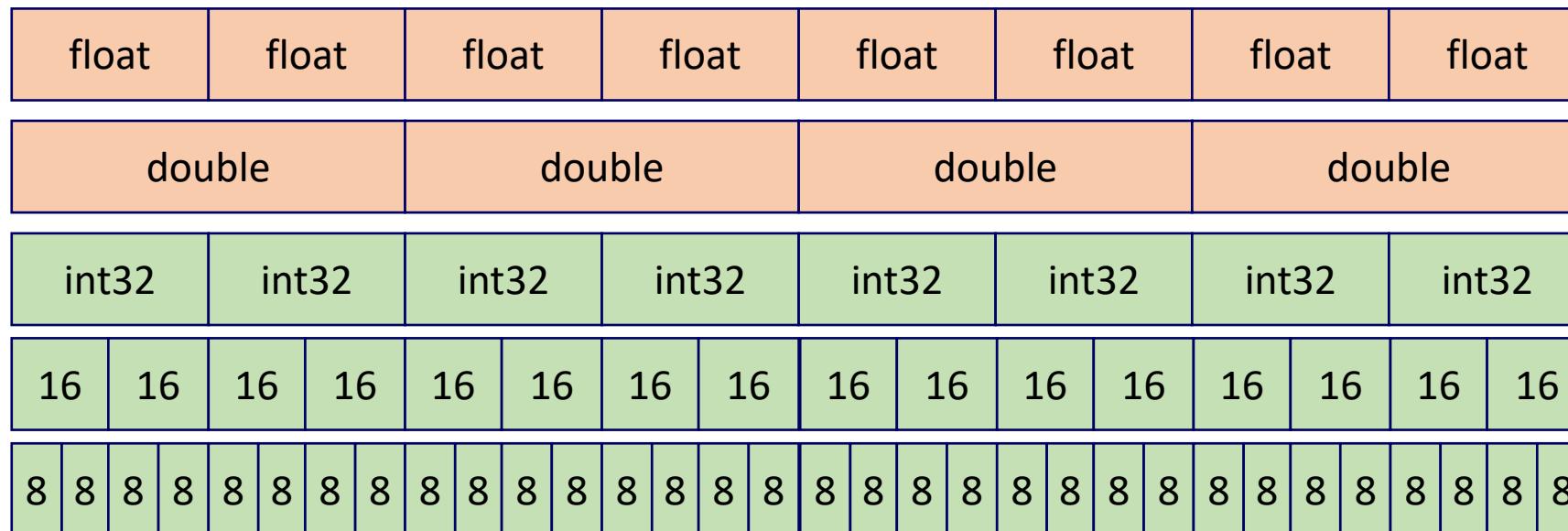
- ❑ XMM0 – XMM15
 - 128-bit registers
 - SSE
- ❑ YMM0 – YMM15
 - 256-bit registers
 - AVX, AVX2
- ❑ ZMM0 – ZMM31
 - 512-bit registers
 - AVX-512

SSE/AVX Data Types

255

0

YMM0

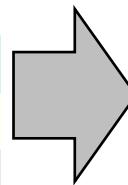


Operation on
32 8-bit values
in one instruction!

Compiler Automatic Vectorization

- ❑ In gcc, flags “-O3 -mavx -mavx2” attempts automatic vectorization
- ❑ Works pretty well for simple loops

```
int a[256], b[256], c[256];
void foo () {
    for (int i=0; i<256; i++) a[i] = b[i] * c[i];
}
```



.L2:

```
    vmovdqa xmm1, XMMWORD PTR b[rax]
    add    rax, 16
    vpmulld xmm0, xmm1, XMMWORD PTR c[rax-16]
    vmovaps XMMWORD PTR a[rax-16], xmm0
    cmp    rax, 1024
    jne    .L2
```

- ❑ But not for anything complex
 - E.g., naïve bubblesort code not parallelized at all

Generated using GCC explorer: <https://gcc.godbolt.org/>

Intel SIMD Intrinsics

- ❑ Use C functions instead of inline assembly to call AVX instructions
- ❑ Compiler manages registers, etc
- ❑ Intel Intrinsics Guide
 - <https://software.intel.com/sites/landingpage/IntrinsicsGuide>
 - One of my most-visited pages...

e.g.,

```
_m256 a, b, c;
```

```
_m256 d = _mm256_fmadd_ps(a, b, c); // d[i] = a[i]*b[i]+c[i] for i = 0 ...7
```

Intrinsic Naming Convention

❑ `_mm<width>_[function]_[type]`

- E.g., `_mm256_fmadd_ps` :
perform fmadd (floating point multiply-add) on
256 bits of
packed single-precision floating point values (8 of them)

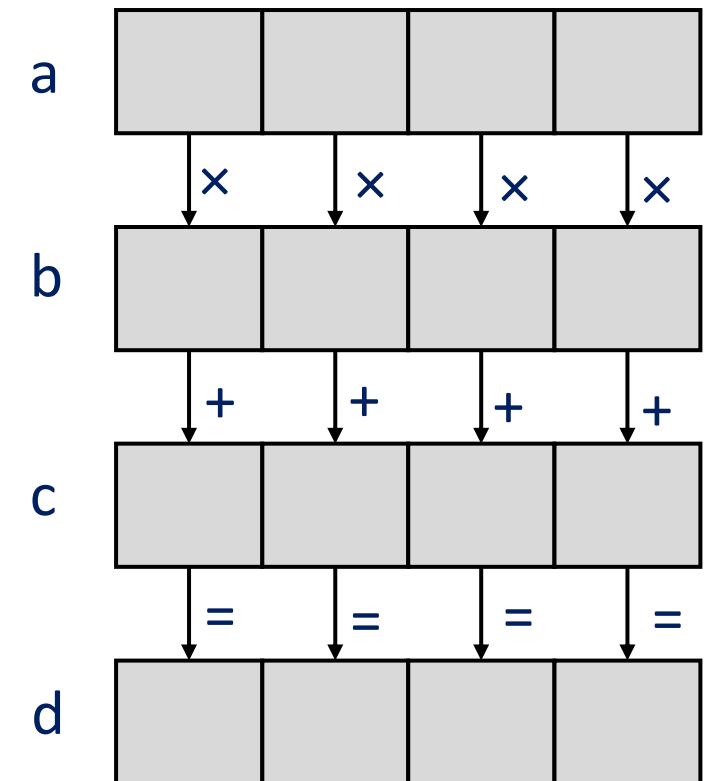
| Width | Prefix |
|-------|----------------------|
| 128 | <code>_mm_</code> |
| 256 | <code>_mm256_</code> |
| 512 | <code>_mm512_</code> |

| Type | Postfix |
|-------------------------|---|
| Single precision | <code>_ps</code> |
| Double precision | <code>_pd</code> |
| Packed signed integer | <code>_epiNNN</code> (e.g., <code>epi256</code>) |
| Packed unsigned integer | <code>_epuNNN</code> (e.g., <code>epu256</code>) |
| Scalar integer | <code>_siNNN</code> (e.g., <code>si256</code>) |

Not all permutations exist! Check guide

Example: Vertical Vector Instructions

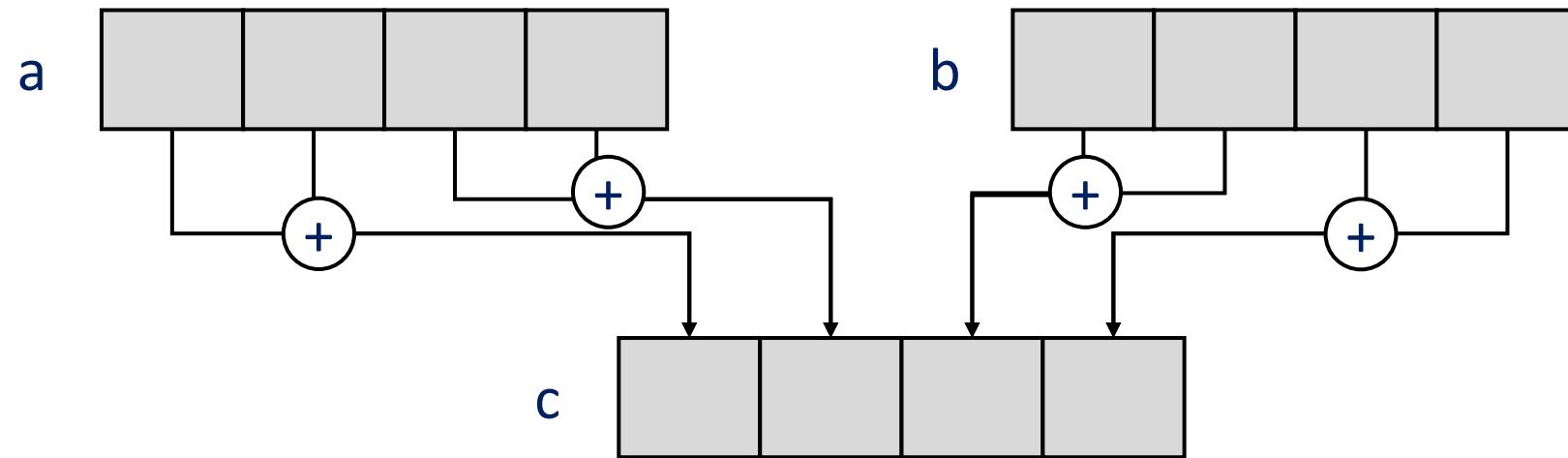
- Add/Subtract/Multiply
 - `_mm256_add/sub/mul/div_ps/pd/epi`
 - Mul only supported for epi32/epu32/ps/pd
 - Div only supported for ps/pd
 - Consult the guide!
- Max/Min/GreaterThan/Equals
- Sqrt, Reciprocal, Shift, etc...
- FMA (Fused Multiply-Add)
 - $(a*b)+c$, $-(a*b)-c$, $-(a*b)+c$, and other permutations!
 - Consult the guide!
- ...



```
_m256 a, b, c;  
_m256 d = _mm256_fmadd_pd(a, b, c);
```

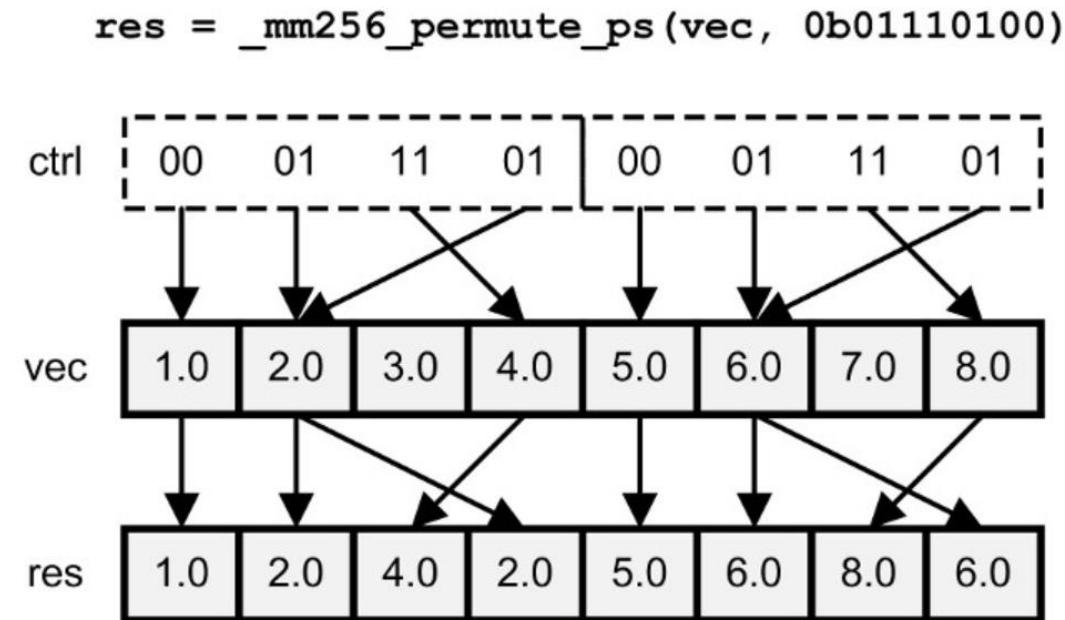
Horizontal Vector Instructions

- ❑ Horizontal add/subtraction
 - Adds adjacent pairs of values
 - E.g., `__m256d _mm256_hadd_pd (__m256d a, __m256d b)`



Shuffling/Permutation

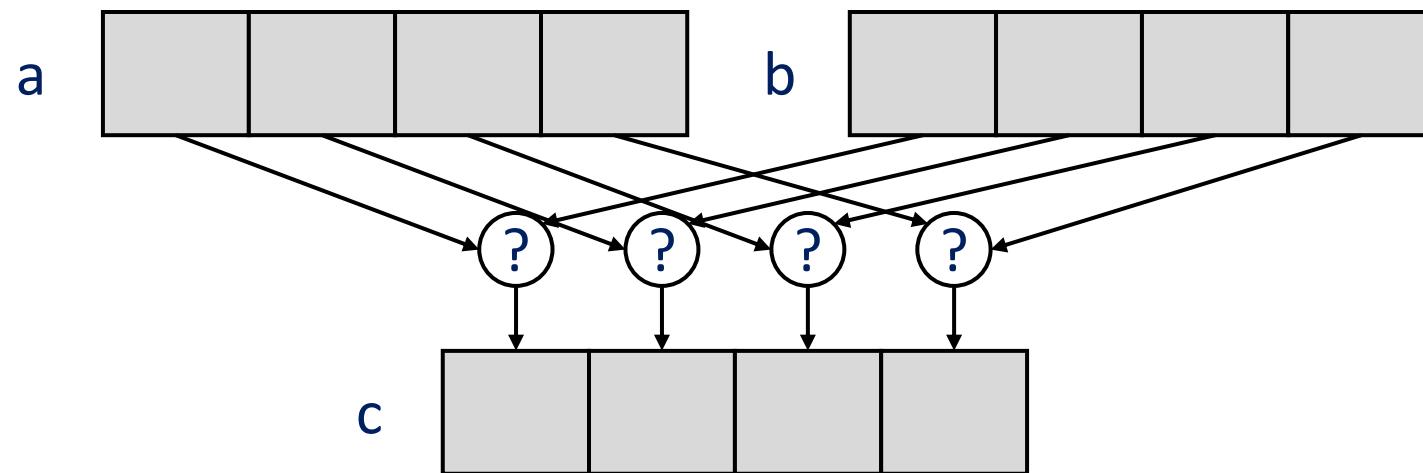
- ❑ Within 128-bit lanes
 - `_mm256_shuffle_ps/pd/...` (a,b, imm8)
 - `_mm256_permute_ps/pd`
 - `_mm256_permutevar_ps/...`
- ❑ Across 128-bit lanes
 - `_mm256_permute2x128/4x64` : Uses 8 bit control
 - `_mm256_permutevar8x32/...` : Uses 256 bit control
- ❑ Not all type permutations exist for each type, but variables can be cast back and forth between types



Matt Scarpino, "Crunching Numbers with AVX and AVX2," 2016

Blend

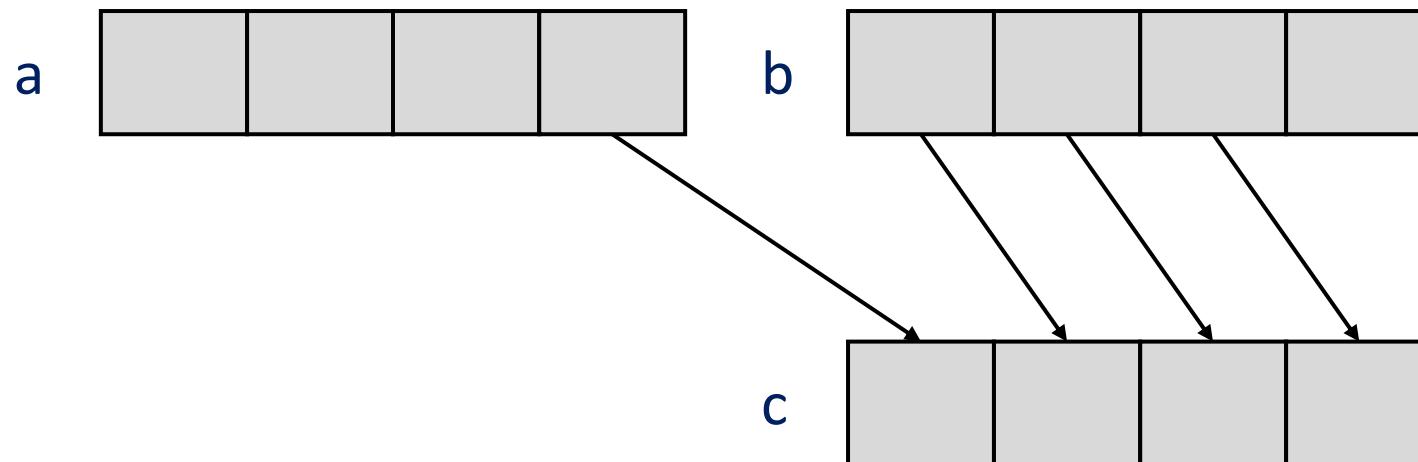
- ❑ Merges two vectors using a control
 - `_mm256_blend_...` : Uses 8 bit control
 - e.g., `_mm256_blend_epi32`
 - `_mm256_blendv_...` : Uses 256 bit control
 - e.g., `_mm256_blendv_epi8`



Alignr

- ❑ Right-shifts concatenated value of two registers, by byte
 - Often used to implement circular shift by using two same register inputs
 - `_mm256_alignr_epi8(a, b, count)`

Example of 64-bit values being shifted by 8

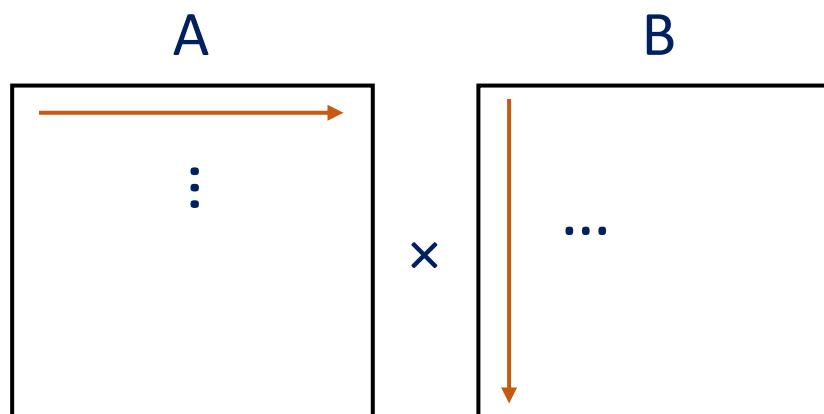


Helper Instructions

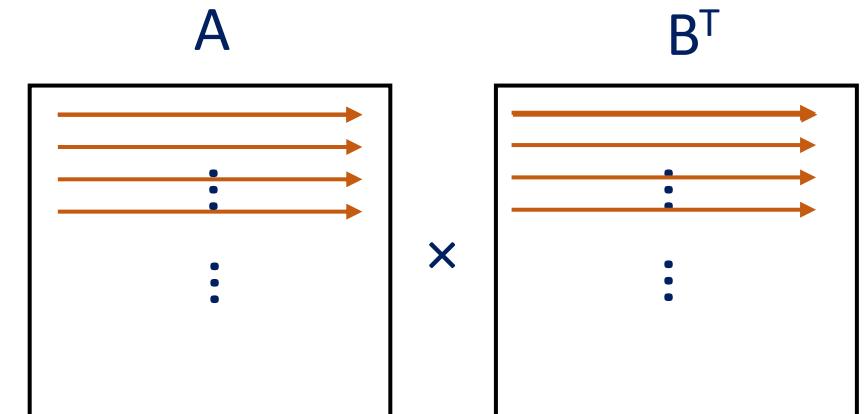
- Cast
 - `_mm256i <-> _mm256`, etc...
 - Syntactic sugar -- does not spend cycles
- Convert
 - 4 floats \leftrightarrow 4 doubles, etc...
- Movemask
 - `_mm256` mask to \rightarrow int imm8
- And many more...

Case Study: Matrix Multiplication

- ❑ Remember simply transposing matrix B brought 6x performance
 - At that point, we are bottlenecked by single-thread processing performance
 - Adding SIMD gets us more!
 - After this we are again bottlenecked by memory, but that is for another time



VS



63.19 seconds

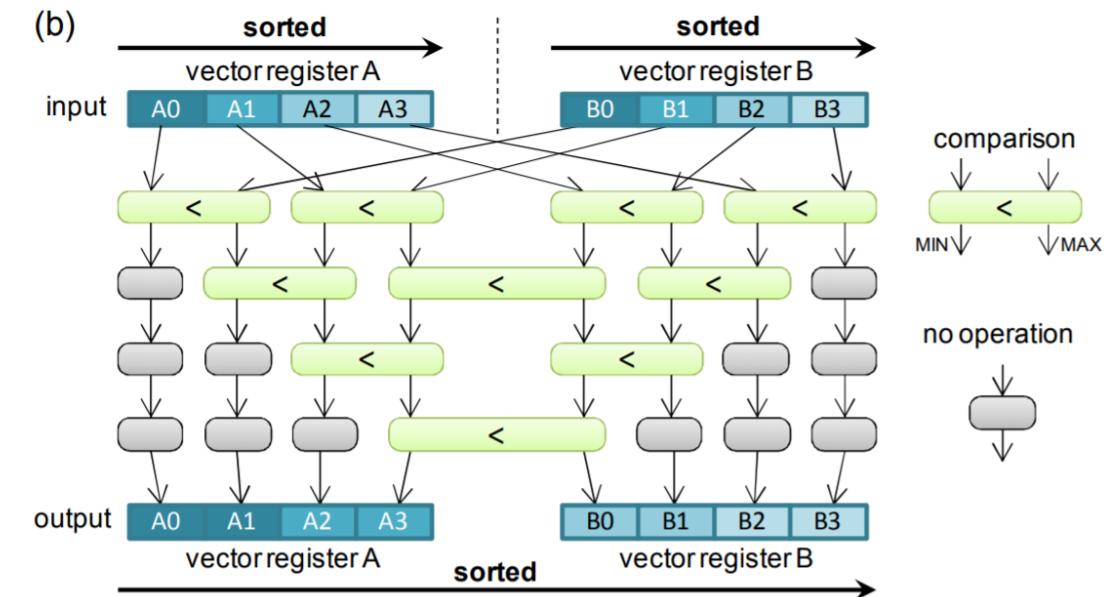
10.69 seconds
(6x performance!)

Case Study: Sorting

- ❑ Important, fundamental application!
- ❑ Can be parallelized via divide-and-conquer
- ❑ How can SIMD help?

The Two Register Merge

- Sort units of two pre-sorted registers, K elements
 - $\text{minv} = A$, $\text{maxv} = B$
 - // Repeat K times
 - $\text{minv} = \min(\text{minv}, \text{maxv})$
 - $\text{maxv} = \max(\text{minv}, \text{maxv})$
 - // circular shift one value down
 - $\text{minv} = \text{alignr}(\text{minv}, \text{minv}, \text{sizeof(int)})$



SIMD And Merge Sort

- ❑ Hierarchically merged sorted subsections
- ❑ Using the SIMD merger for sorting
 - vector_merge is the two-register sorter from before

```
aPos = bPos = outPos = 0;
vMin = va[aPos++];
vMax = vb[bPos++];
while (aPos < aEnd && bPos < bEnd) {
    /* merge vMin and vMax */
    vector_merge(vMin, vMax);

    /* store the smaller vector as output*/
    vMergedArray[outPos++] = vMin;

    /* load next vector and advance pointer */
    /* a[aPos*4] is first element of va[aPos] */
    /* and b[bPos*4] is that of vb[bPos] */
    if (a[aPos*4] < b[bPos*4])
        vMin = va[aPos++];
    else
        vMin = vb[bPos++];
}
```

Inoue et.al., “SIMD- and Cache-Friendly Algorithm for Sorting an Array of Structures,” VLDB 2015

Topic Under Active Research!

- ❑ Papers being written about...
 - Architecture-optimized matrix transposition
 - Register-level sorting algorithm
 - Merge-sort
 - ... and more!
- ❑ Good find can accelerate your application kernel Nx

Processor Microarchitectural Effects on Power Efficiency

- ❑ The majority of power consumption of a CPU is not from the ALU
 - Cache management, data movement, decoding, and other infrastructure
 - Adding a few more ALUs should not impact power consumption
- ❑ Indeed, 4X performance via AVX does not add 4X power consumption
 - From i7 4770K measurements:
 - Idle: 40 W
 - Under load : 117 W
 - Under AVX load : 128 W

